

# Quantum Algorithms for Approximating Partition Functions

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(w/ Arjan Cornelissen)

# Estimating partition functions

Given a **classical** Hamiltonian  $H : \Omega \rightarrow \{0, 1, \dots, n\}$

and an inverse temperature  $\beta$

approximate the partition function  $Z(\beta) = \text{Tr}(e^{-\beta H})$

Gibbs state:  $|\pi_\beta\rangle \propto \sum_i e^{-\beta H(i)} |i\rangle$

Energy basis = standard basis

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Gibbs state:  $|\pi_\beta\rangle = \frac{1}{\sqrt{Z(\beta)}} \sum_i e^{-\beta H(i)} |i\rangle$

# Quantum algorithms hold the promise of faster **statistical estimation**

- Heisenberg vs shot-noise limits, quantum metrology
- Quantum **phase estimation**
- Quantum counting, quantum mean estimation
- ...

## How far can we push this advantage?

- **Generic** methods to turn classical stat. algos into (faster) quantum algos?
- **Quadratic** speedups?
- Highly **structured** estimation tasks?

# Estimating partition functions

- Partition functions are **ubiquitous**
  - statistical physics
  - combinatorics (counting matchings, independent sets...)
  - linear algebra (permanents)
  - convex geometry (volume of a body)
  - machine learning (graphical models)
  - ...

# Estimating partition functions

- Partition functions are ubiquitous
- Related to **counting** (generating functions) and **phase transitions**
- **Exact** computation is often **#P-hard**
- ... but there exists efficient **approximation** methods (**MCMC**, Taylor's approximation, correlation decay, ...)

# Main Result

The classical Markov Chain Monte Carlo (MCMC) algorithms★ for estimating  $Z(\beta)$  can be speed-up by quantum algorithms running in time

$$\approx \log^{1/4} |\Omega| \times \sqrt{\text{time of classical MCMC algo.}}$$

**This class of methods includes the best known algorithms for:**

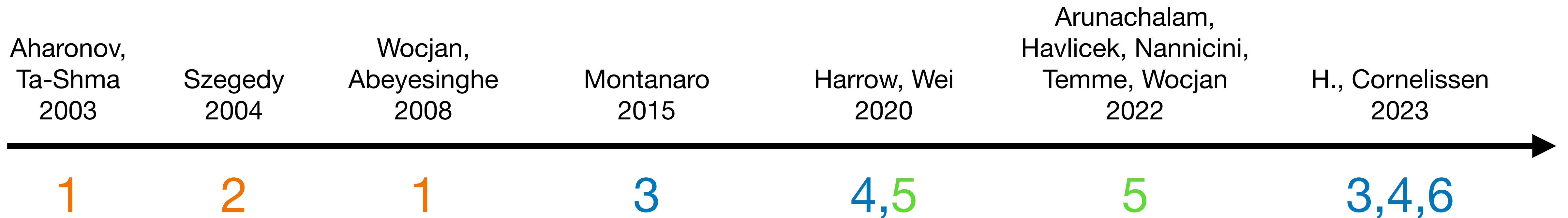
- number of independent sets, colorings, matchings
- Ising, Potts, monomer-dimer,... models
- volume of convex bodies
- permanent of nonnegative matrices

$$\geq \log |\Omega|$$

( $\Omega$  = configuration space)

# Main Result

$$\approx \log^{1/4} |\Omega| \times \sqrt{\text{time of classical MCMC algo.}}$$



1. Quantum simulated annealing

2. Szegedy's quantum walk

3. Variance reduction

4. Non-destructive estimation

5. Cooling schedule computation

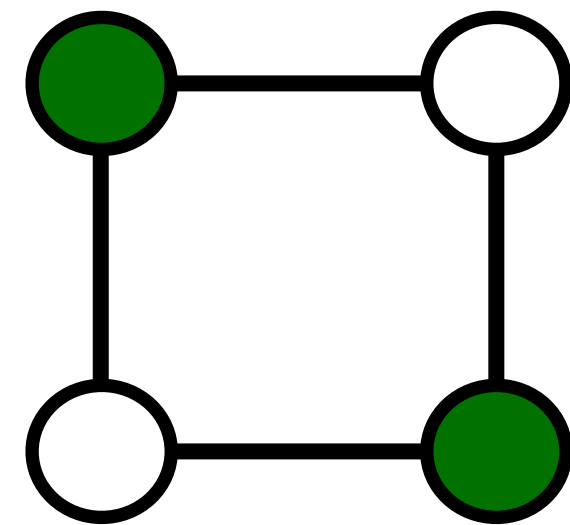
6. Unbiased estimation



# The independent set partition function

# Independent set

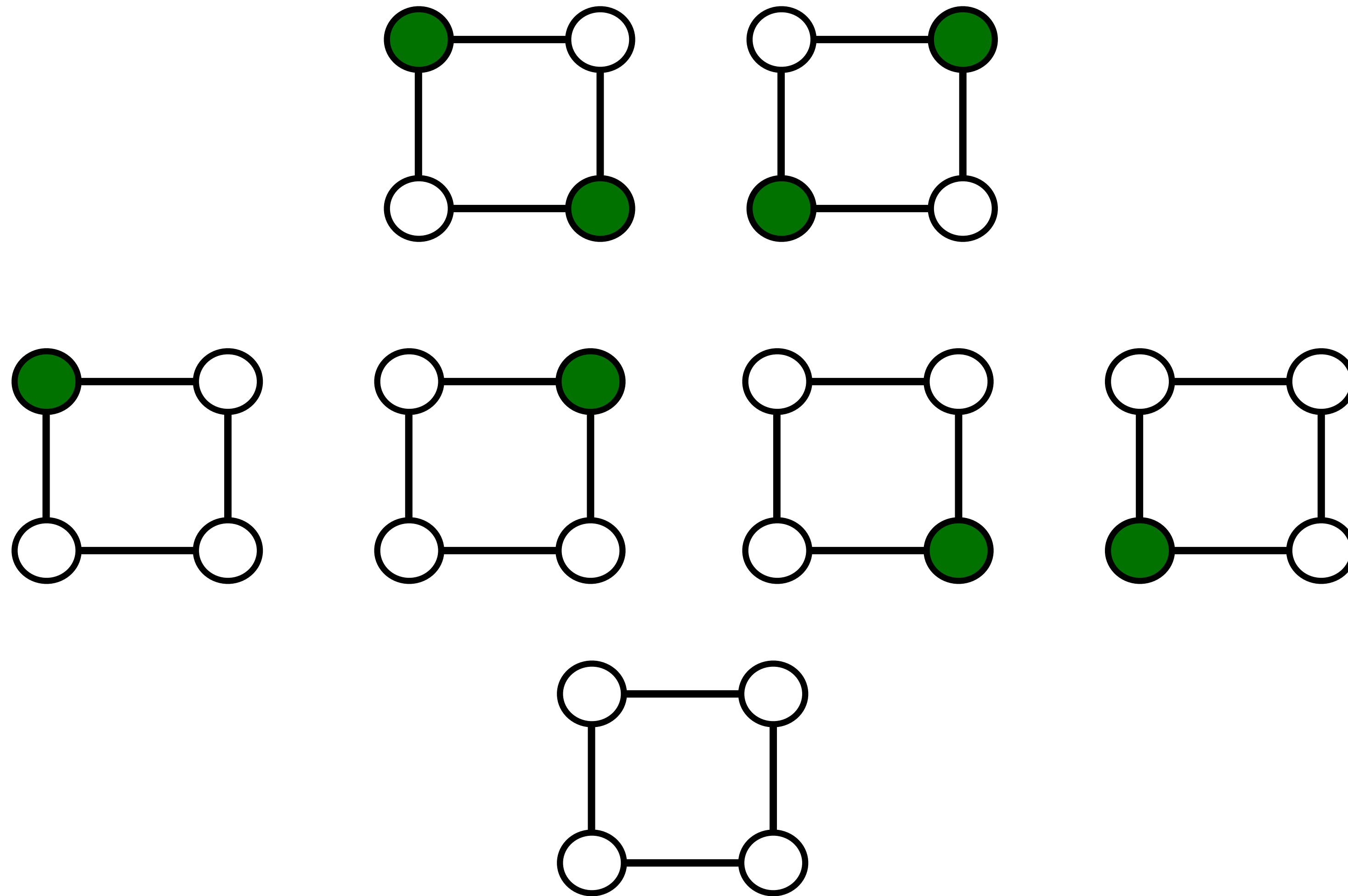
= subset of non-adjacent vertices



● = occupied

Hard-core gas model in statistical physics

# independent sets = 7



Input: graph  $G$

Output: # independent sets of  $G$

#P-hard in many regimes

Bipartite graphs

[Provan, Ball'83]

3-regular graphs

[Dyer, Greenhill'00]

...

Exact counting



Approximate counting?

Input: graph  $G$  and  $\epsilon \in (0,1)$

Output:  $S$  s.t.  $(1 - \epsilon) \#ind \leq S \leq (1 + \epsilon) \#ind$

- Graphs with maximum vertex-degree  $\leq 5$ :  $(n = \#vertices)$

Classical algorithms

$$\tilde{O}(n^2/\epsilon^2)$$

[Štefankovič, Vempala, Vigoda'09]

[Chen, Liu, Vigoda'21]

Quantum algorithms

$$\tilde{O}(n^{5/4}/\epsilon)$$

[H. Cornelissen'23]

- Graphs with maximum vertex-degree  $> 5$ :

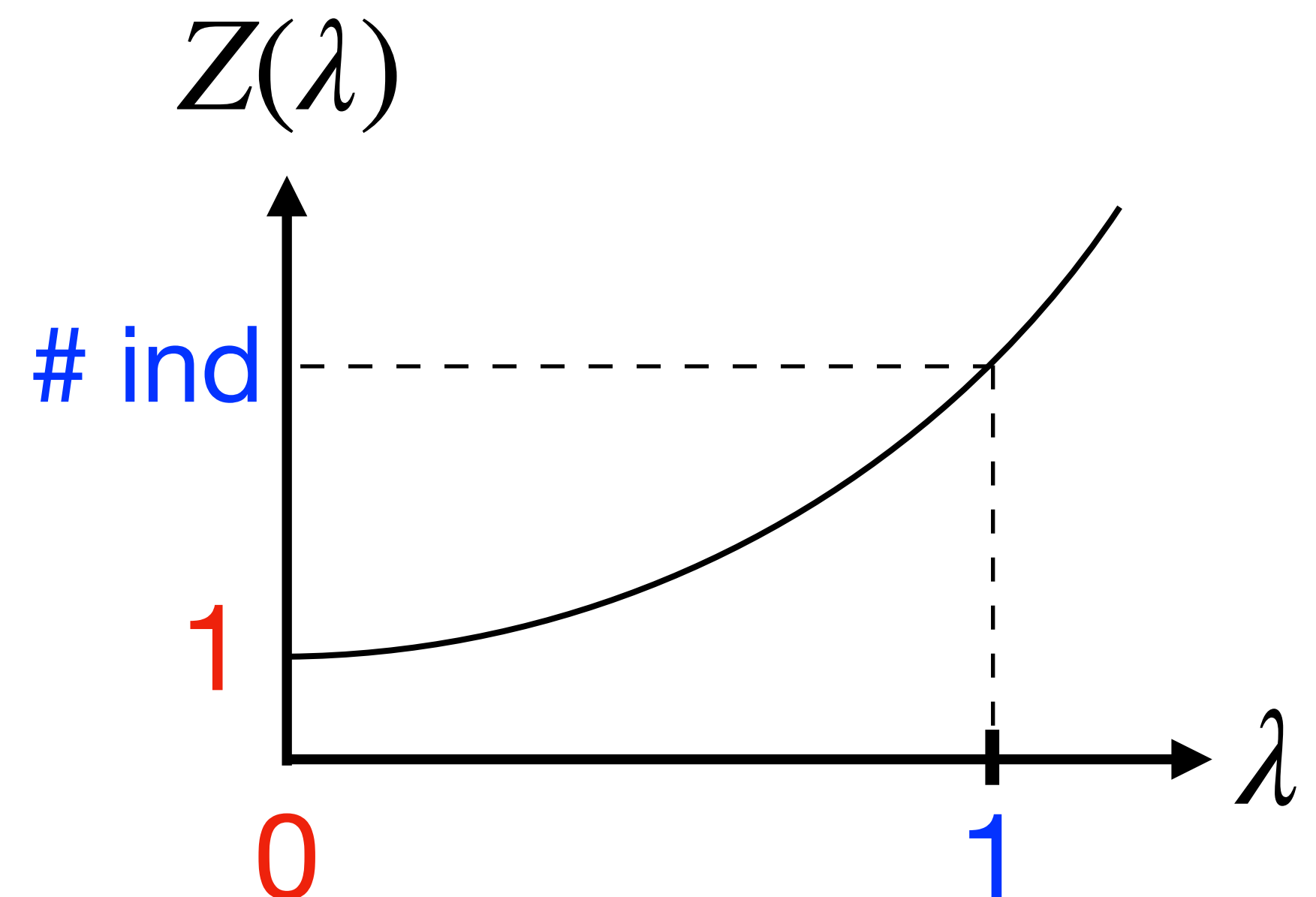
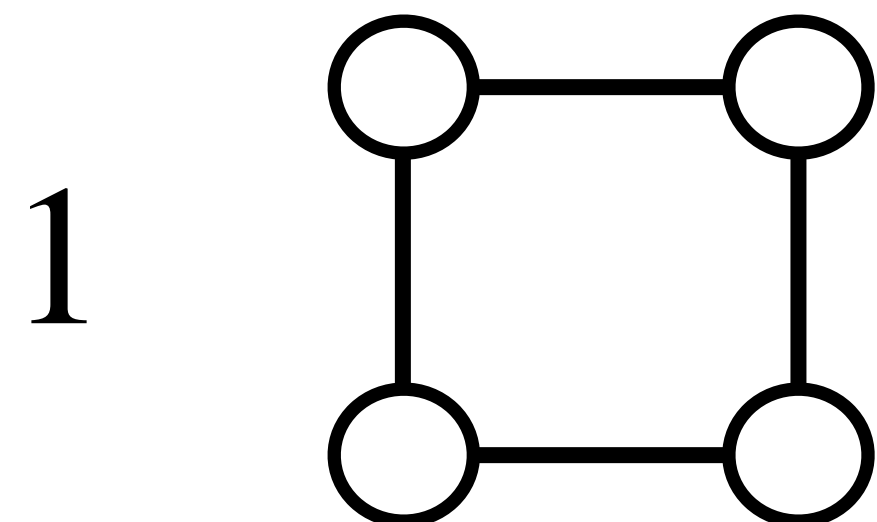
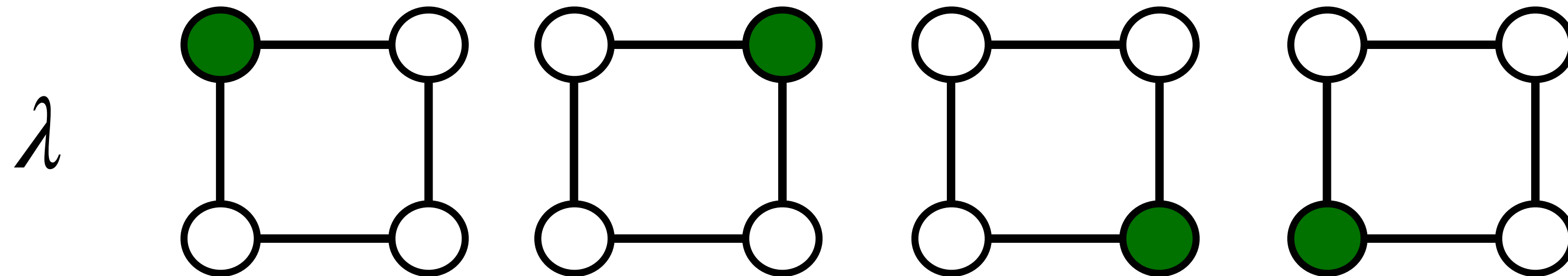
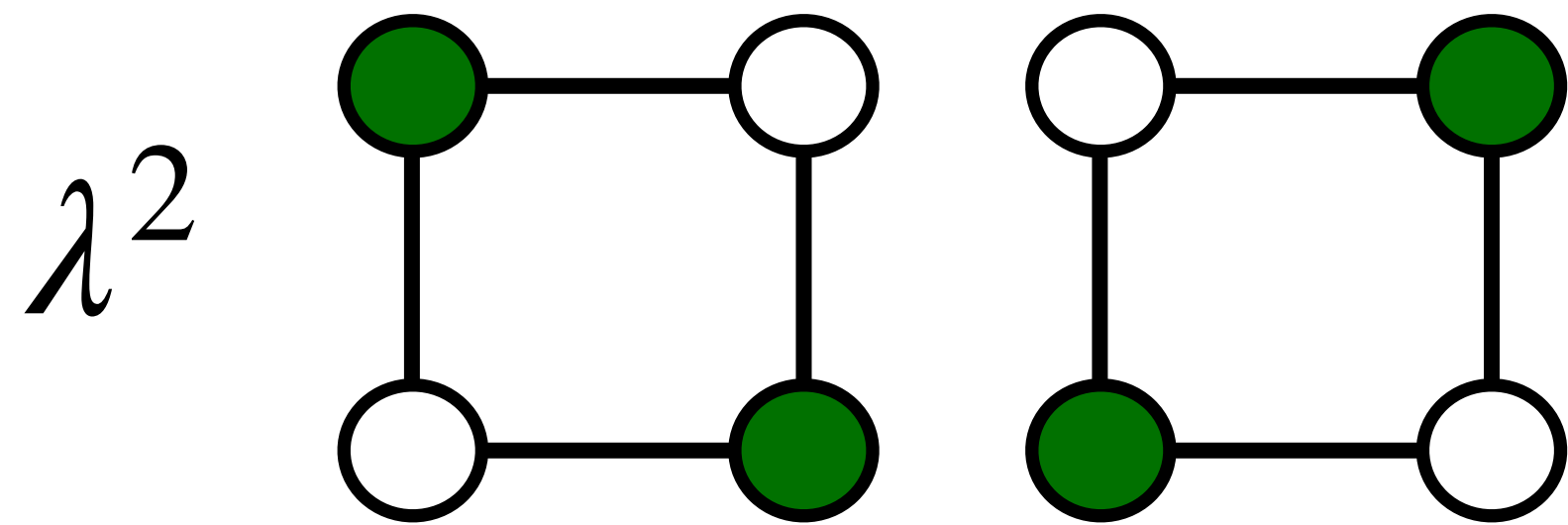
No FPRAS unless  $NP = RP$  [Sly'10]

# Classical MCMC method for independent sets

# Weighted independent sets

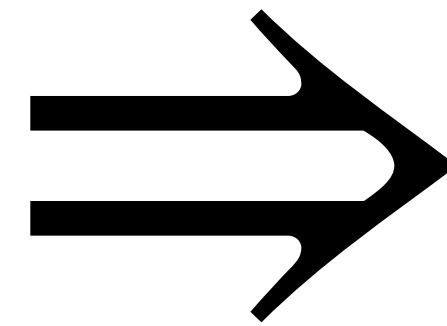
$\lambda = \text{fugacity}$

Partition function:  $Z(\lambda) = \sum_{I \text{ ind. set}} \lambda^{|I|}$



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Sampling



Estimating

Independent set  $I$

with probability  $\pi_\lambda(I) = \frac{\lambda^{|I|}}{Z(\lambda)}$

Gibbs sampling

Ratio estimation

$$\frac{Z(\lambda')}{Z(\lambda)} = \mathbb{E}_{I \sim \pi_\lambda} \left( \frac{\lambda'}{\lambda} \right)^{|I|}$$

Unbiased estimator



# Gibbs sampling

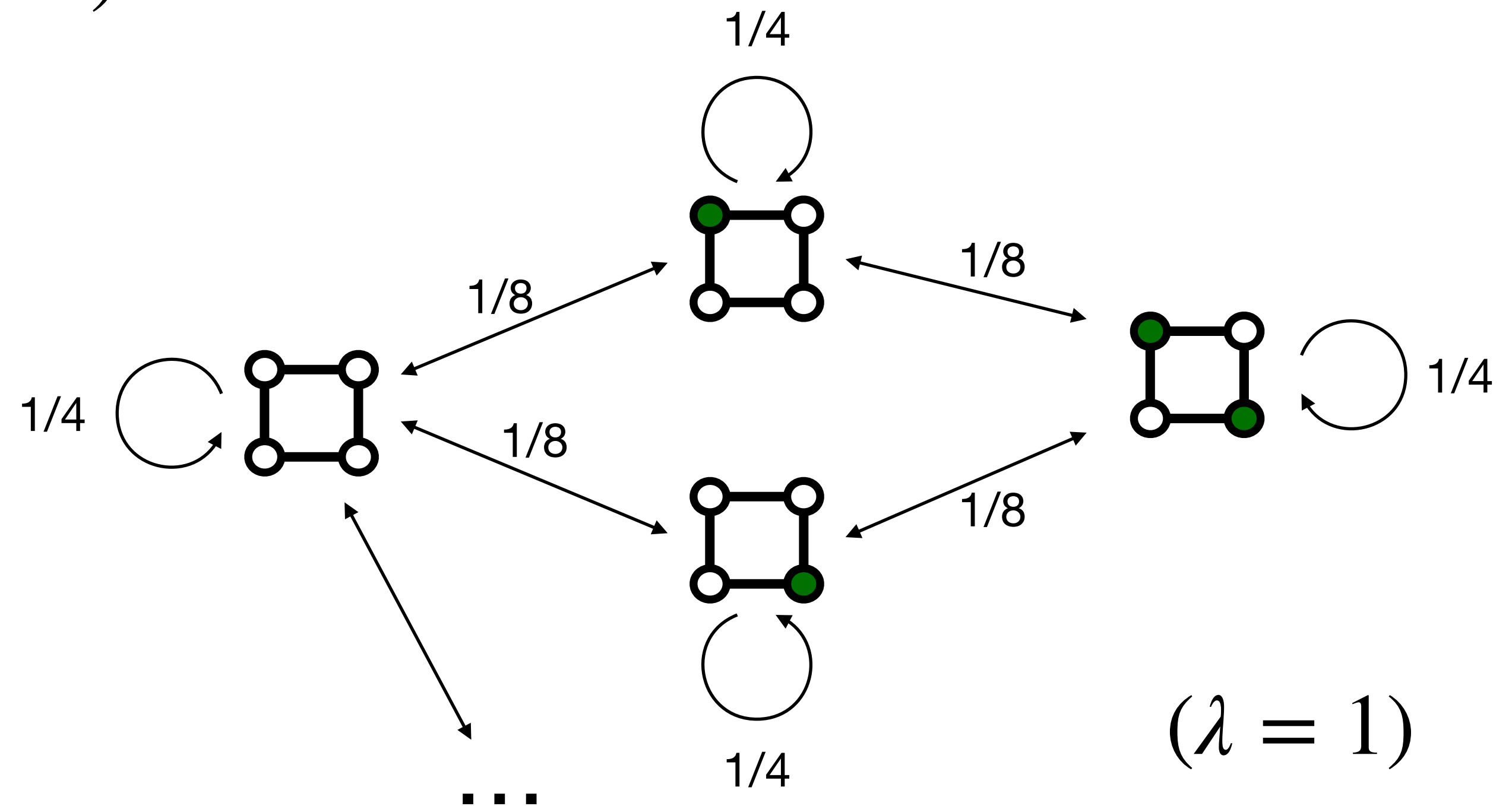
Run a **Markov Chain** with **stationary distribution**  $\pi_\lambda$  :

1. Choose a vertex uniformly at random
2. Make it occupied with proba  $\lambda/(\lambda + 1)$  if the set remains independent

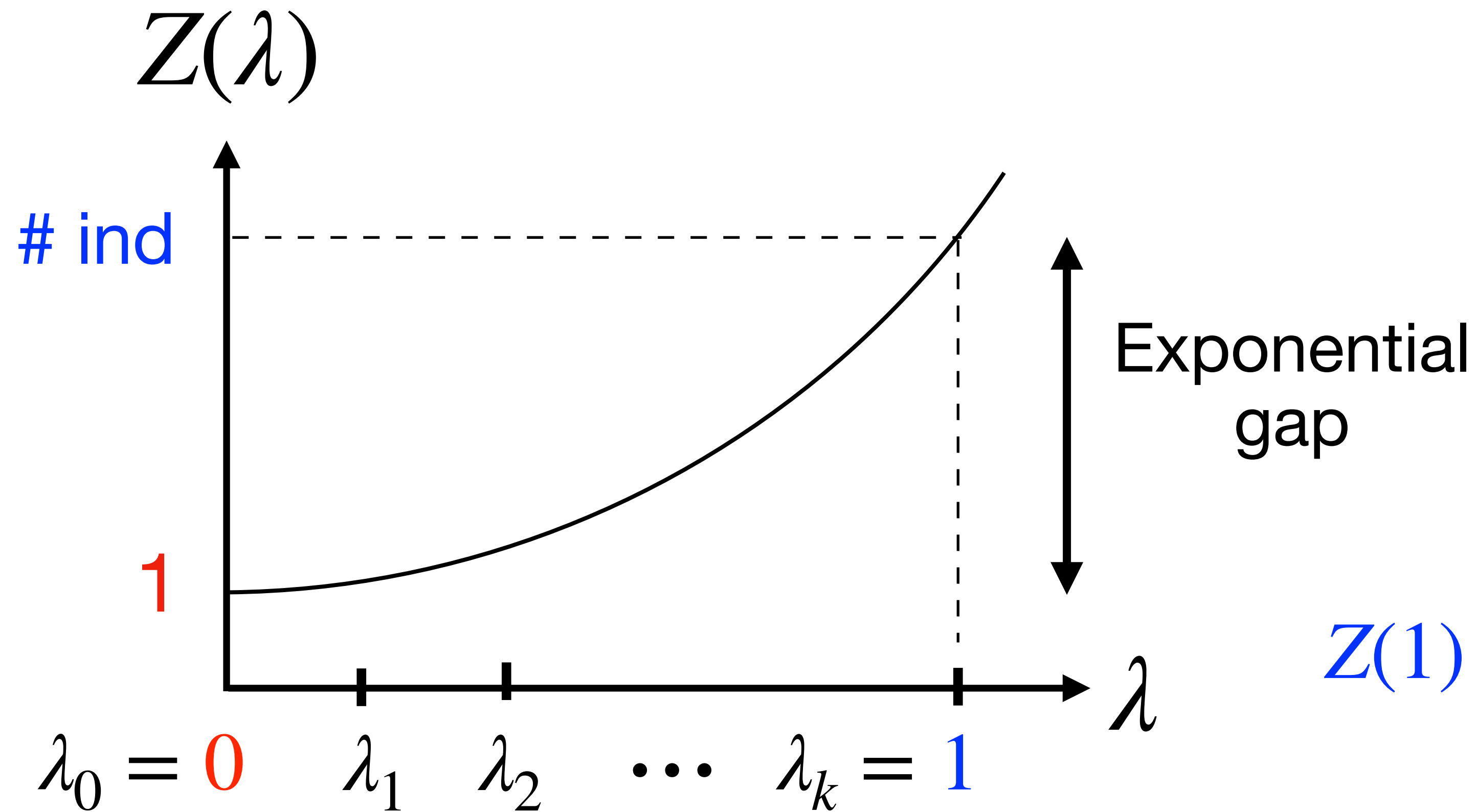
Glauber dynamics

Mixing time =  $O(n \log n)$

[Chen,Liu,Vigoda'21]



# Partition function estimation



Slowly move from an estimate of  $Z(0)$  to an estimate of  $Z(1)$

$$Z(1) = \frac{Z(\lambda_1)}{Z(\lambda_0)} \cdot \frac{Z(\lambda_2)}{Z(\lambda_1)} \cdot \dots \cdot \frac{Z(\lambda_k)}{Z(\lambda_{k-1})}$$

$$\begin{array}{ccc} \uparrow & \uparrow & \uparrow \\ E_{\pi_{\lambda_0}} \left( \frac{\lambda_1}{\lambda_0} \right)^{|I|} & E_{\pi_{\lambda_1}} \left( \frac{\lambda_2}{\lambda_1} \right)^{|I|} & E_{\pi_{\lambda_{k-1}}} \left( \frac{\lambda_k}{\lambda_{k-1}} \right)^{|I|} \end{array}$$

Balance **length** of the sequence and **variance** of each estimator

# Quantum MCMC method for independent sets

## Two interlaced branches of work:

1

Converting classical reversible **Markov chains** into **quantum walks**

2

Estimating **expectation values** on the 1-eigenvector of a quantum walk

## Glauber dynamics

Sample ind. set  $I \sim \pi_\lambda$  in time  $O(n \log n)$

No speedup for **sampling**? Diameter(Markov Chain) =  $n$

We can simulate in time  $\tilde{O}(\sqrt{n})$  the **reflection** through the state:

$$|\pi_\lambda\rangle = \sum_I \sqrt{\pi_\lambda(I)} |I\rangle$$

(quantum Gibbs sample)

Classical  
sampling resource

$$I \sim \pi_\lambda$$

Time  $\tilde{O}(n)$

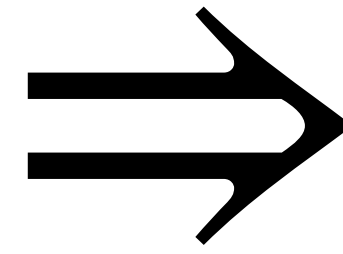
Quantum  
sampling resource

$$2 |\pi_\lambda\rangle\langle\pi_\lambda| - I$$

Time  $\tilde{O}(\sqrt{n})$

# Szegeedy Quantum Walk

Markov chain (ergodic reversible)  
with stationary distribution  $\pi_\lambda$



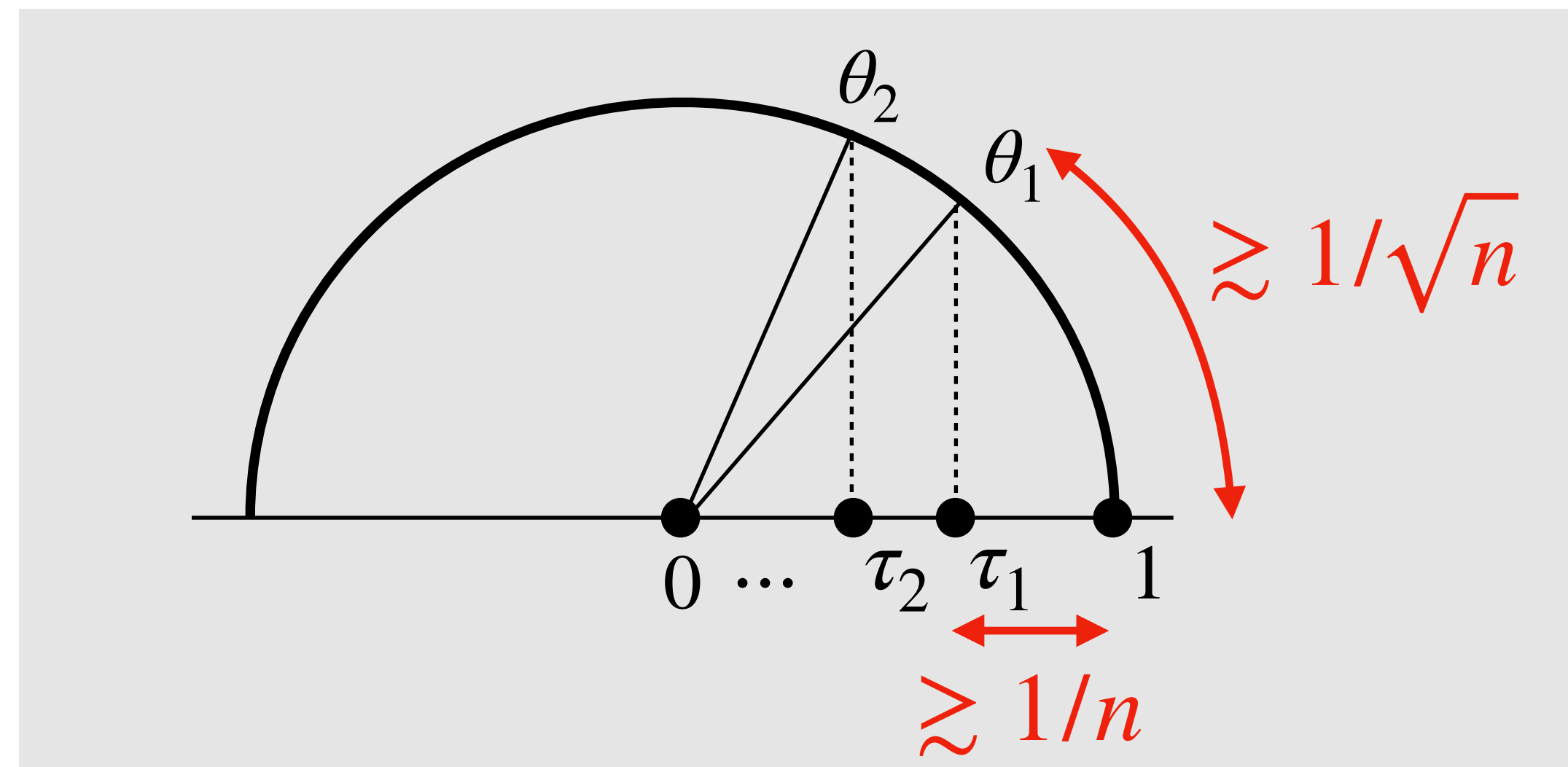
Quantum walk with  
1-eigenvector  $|\pi_\lambda\rangle$

Transition matrix  $P$

Unitary  $W(P)$

$$\text{Spec}(P) = \{1, \tau_1, \tau_2, \dots\}$$

$$\text{Spec}(W(P)) = \{1, e^{\pm 2i\theta_1}, e^{\pm 2i\theta_2}, \dots\}$$



Run  $\tilde{O}(\sqrt{n})$  steps of **quantum phase estimation** on the quantum walk operator  $W(P)$  to simulate the reflection

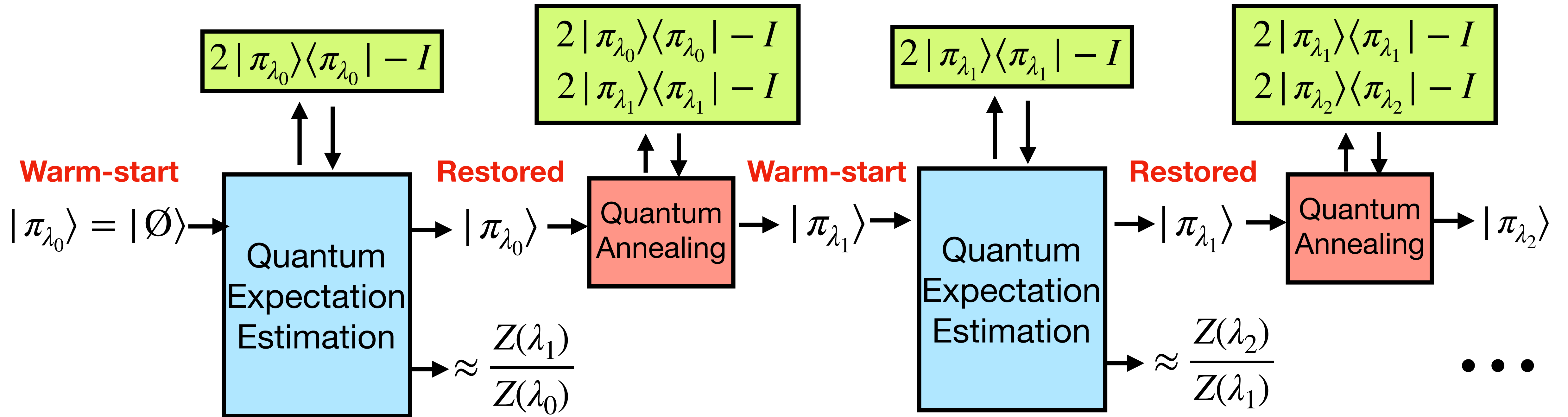
$$2|\pi_\lambda\rangle\langle\pi_\lambda| - I$$

How to use this reflection to estimate **expectation values** on  $|\pi_\lambda\rangle$  ?

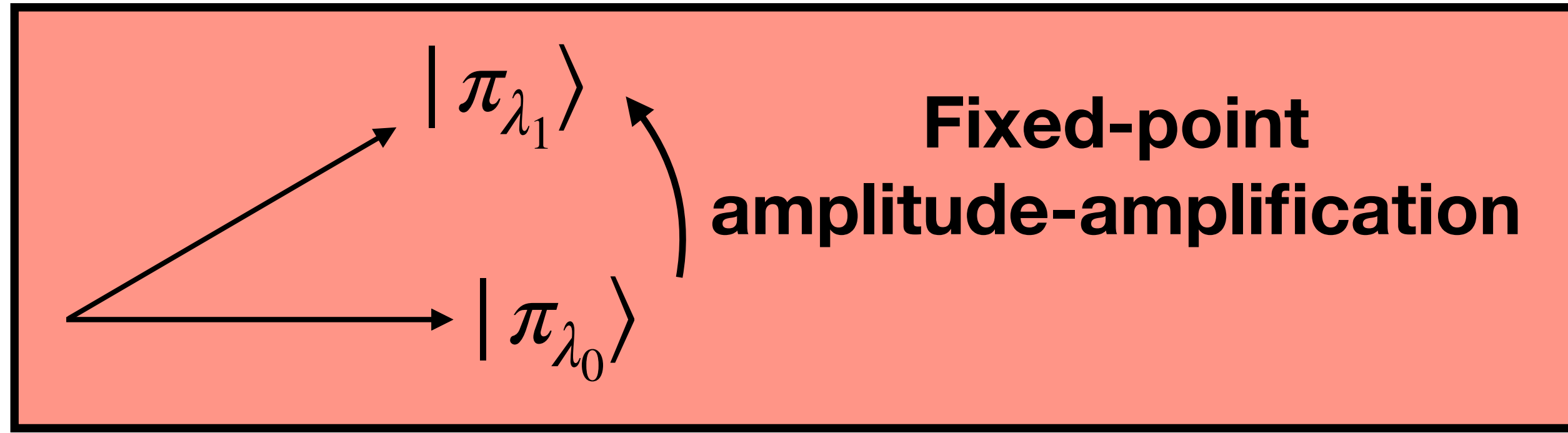
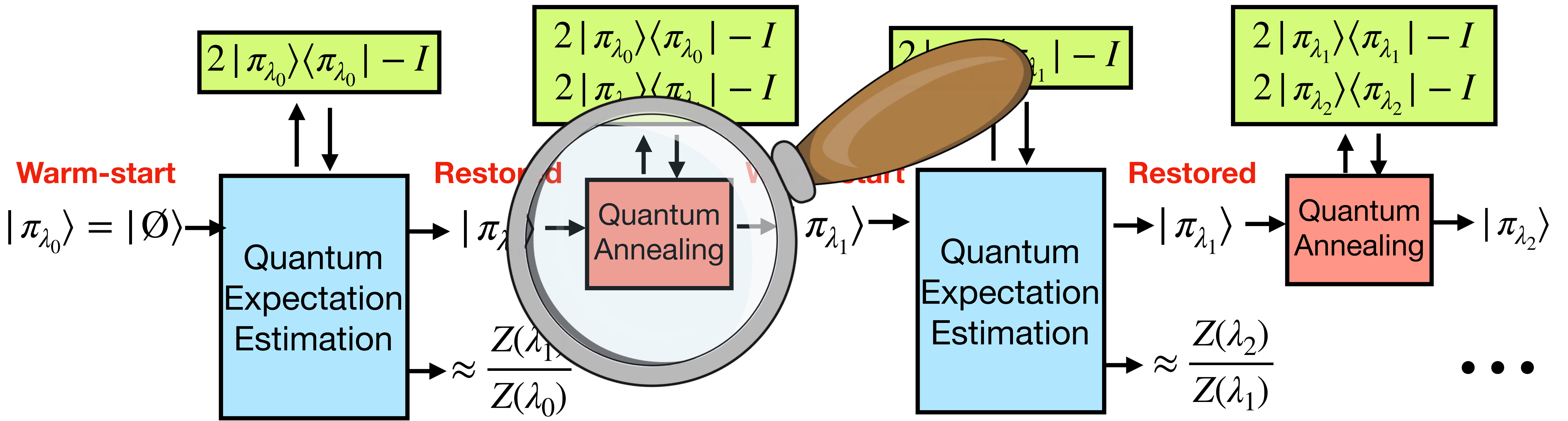
*Example:*  $\frac{Z(\lambda')}{Z(\lambda)} = \langle\pi_\lambda| e^{-(\lambda'-\lambda)H} |\pi_\lambda\rangle$



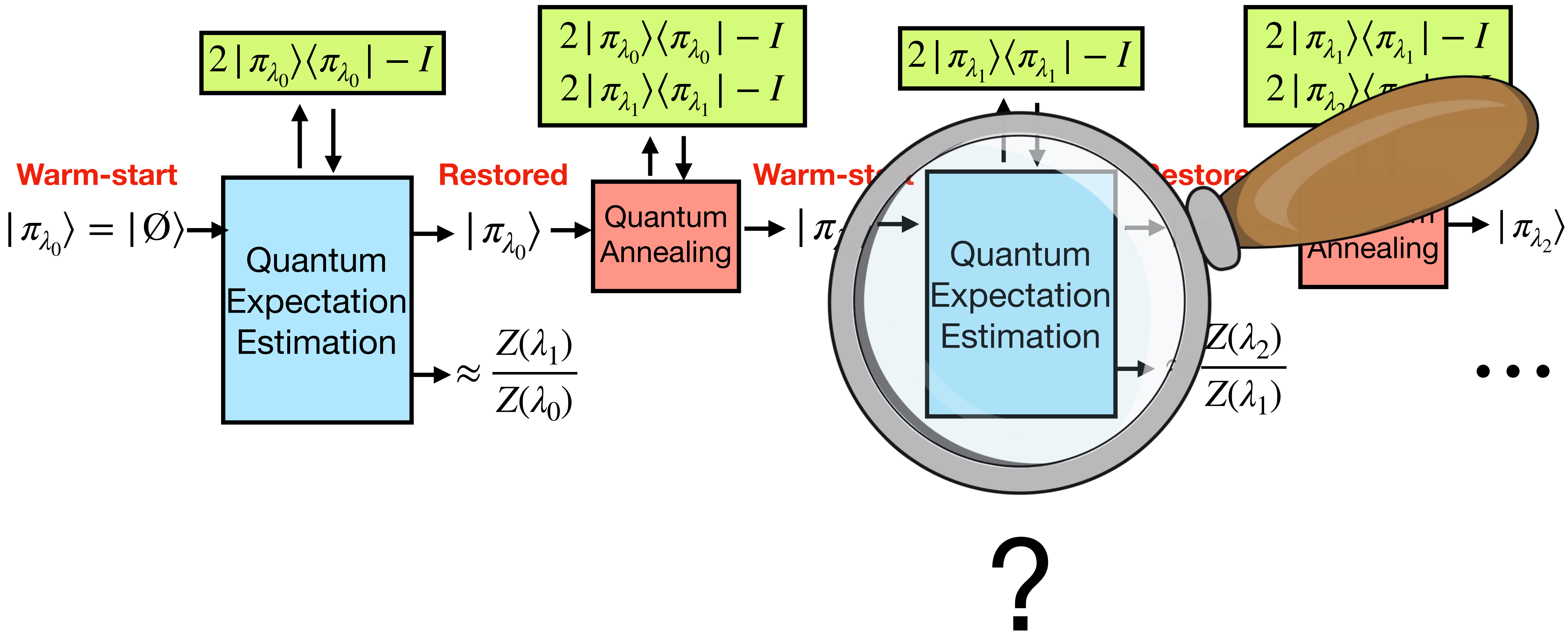
# Quantum partition estimation framework



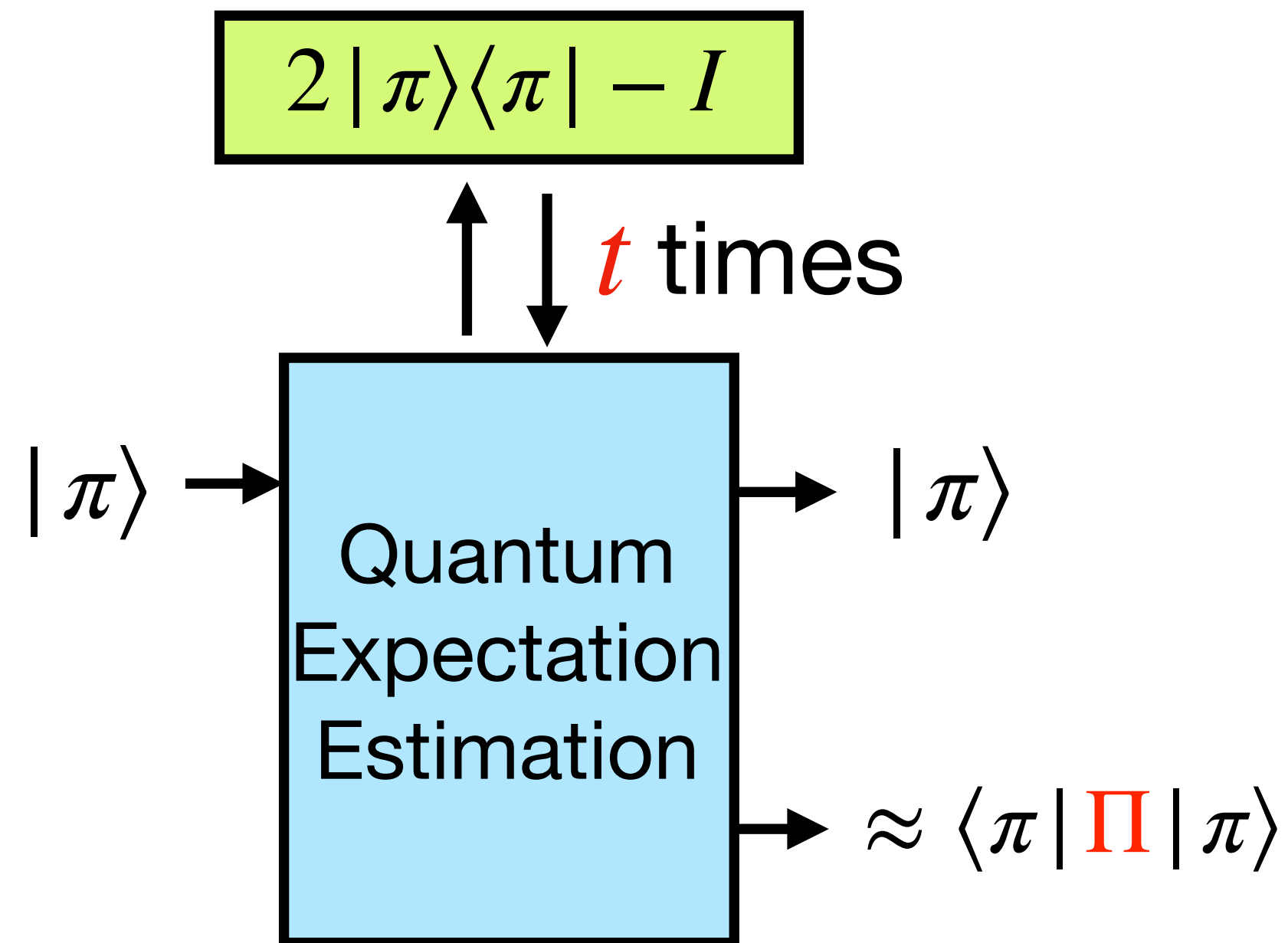
# Quantum partition estimation framework



# Quantum partition estimation framework



# Base case: projection observables



**Input:**

- description of projector  $\Pi$
- 1 copy of  $|\pi\rangle$
- $t$  access to  $2|\pi\rangle\langle\pi| - I$   
(morally,  $t = \text{sample complexity}$ )

**Output:**

- 1 new copy of  $|\pi\rangle$
- estimate of  $\langle \pi | \Pi | \pi \rangle$   
(unbiased + low variance)

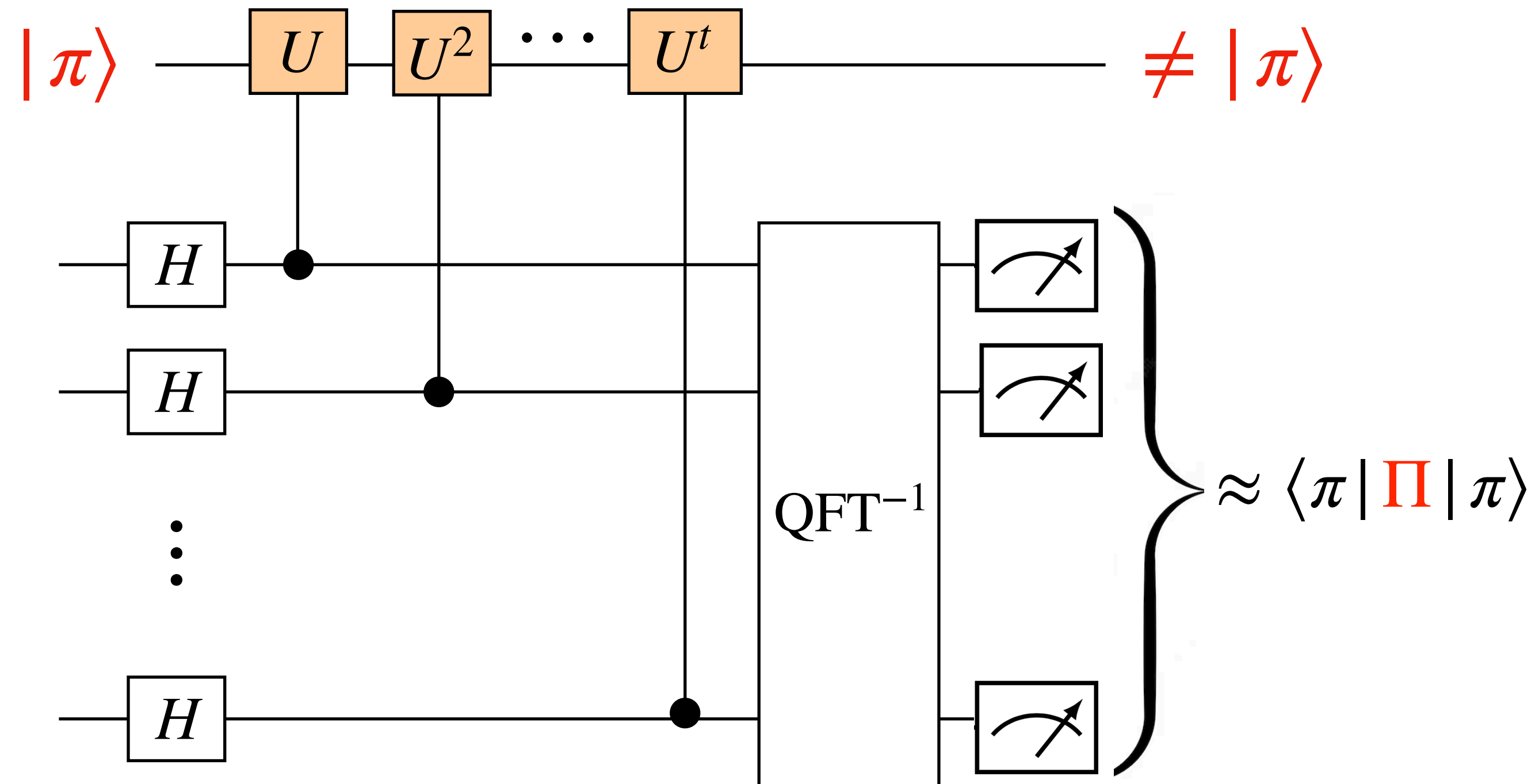
# Base case: projection observables

## Grover operator

$$U = (2|\pi\rangle\langle\pi| - I)(2\Pi - I)$$

$|\pi\rangle$  is a superposition over the  $(\pm 2 \sin^{-1} \sqrt{\langle\pi|\Pi|\pi\rangle})$ -eigenvectors of  $U$

## Quantum Phase estimation

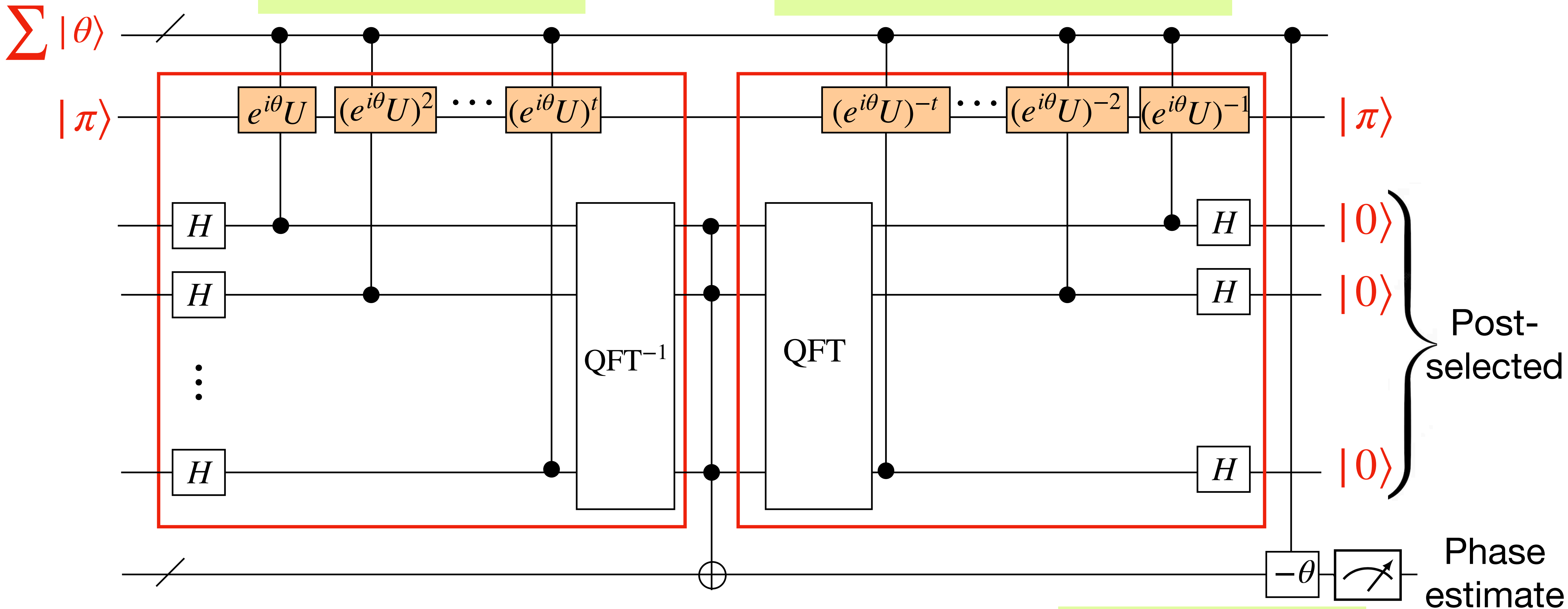


# New phase estimation primitive

Random phase shift

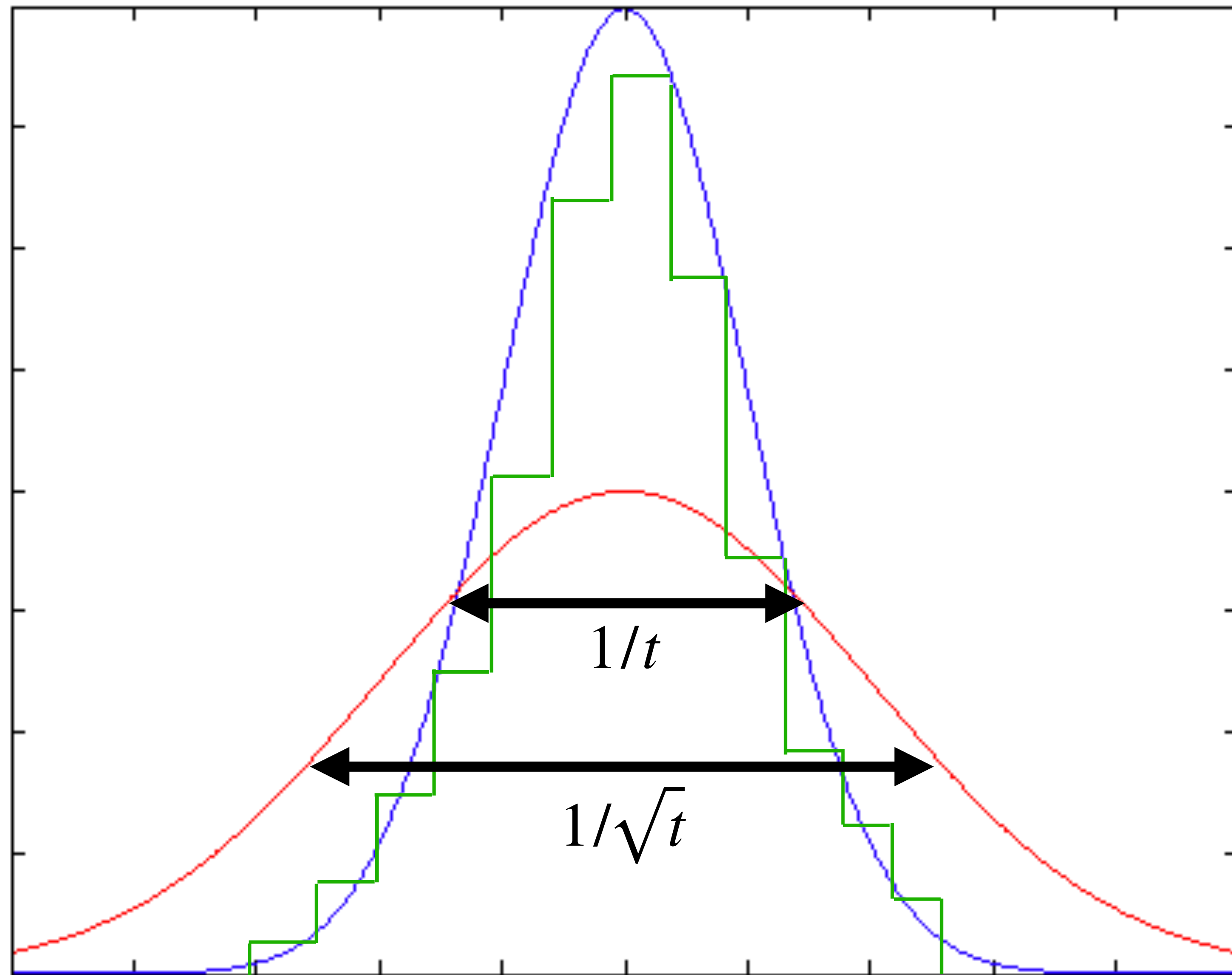
Phase estimation

Inverse Phase estimation



Phase correction

# New phase estimation primitive



Outcome distributions after **t steps** of:

- Empirical mean (classical)
- Standard phase estimation
- Enhanced phase estimation

Open questions



# Classical partition functions

The partition function  $Z(\beta)$  of **classical** Hamiltonians can be estimated in

$$\approx \log^{1/4} |\Omega| \times \sqrt{\text{time of classical MCMC algo.}}$$

- Can we remove the  $\log^{1/4} |\Omega|$  factor to get a full-quadratic speedup?

⇒ counting independent sets in time  $\tilde{O}(n/\epsilon)$  ?

⇒ quantum sampling  $\pi_\beta$  and estimating  $Z(\beta)$  would have  $\approx$  complexity!

- ... or show a **polynomial separation**  $T_{\text{estim.}} \geq T_{\text{sampl.}}^c$  for some Hamiltonian

# Quantum partition functions

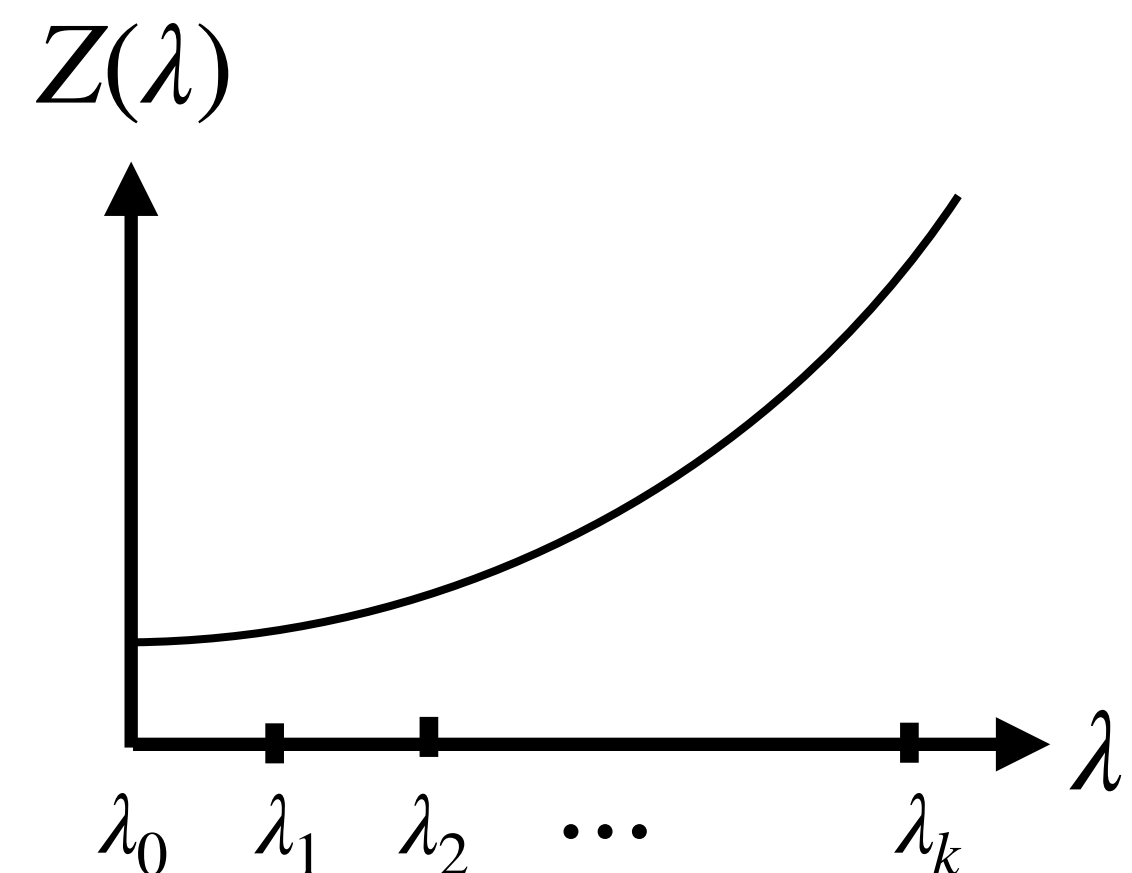
$$\text{Estimate } Z(\beta) = \text{Tr}(e^{-\beta H})$$

when  $H$  is a quantum Hamiltonian

Efficient **reduction** from quantum partition function estimation to quantum Gibbs sampling?

$$\frac{Z(\lambda')}{Z(\lambda)} = \text{Tr}(e^{-(\lambda'-\lambda)H} \pi_\lambda)$$

Ratio estimation?



Cooling schedule?