

# Quantum Chebyshev's Inequality and Applications

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## Mean Estimation Problem

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**In practice:** we often know  $\Delta^2 \geq \frac{E(X^2)}{E(X)^2} = \frac{\text{Var}(X^2)}{E(X)^2} + 1 \rightarrow$  take  $\frac{\Delta^2}{\epsilon^2}$  samples

## Counting with Markov chain Monte Carlo methods:

Counting vs. sampling [Jerrum, Sinclair'96] [Štefankovič et al.'09], Volume of convex bodies [Dyer, Frieze'91], Permanent [Jerrum, Sinclair, Vigoda'04]

## Data stream model:

Frequency moments, Collision probability [Alon, Matias, Szegedy'99] [Monemizadeh, Woodruff'] [Andoni et al.'11] [Crouch et al.'16]

## Testing properties of distributions:

Closeness [Goldreich, Ron'11] [Batu et al.'13] [Chan et al.'14], Conditional independence [Canonne et al.'18]

## Estimating graph parameters:

Number of connected components, Minimum spanning tree weight [Chazelle, Rubinfeld, Trevisan'05], Average distance [Goldreich, Ron'08], Number of triangles [Eden et al. 17]

etc.



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**Question:** can we estimate  $E(X)$  with less samples in the quantum setting?

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$$\frac{B}{\epsilon^2 E(X)}$$

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$$\frac{\Delta^2}{\epsilon^2} \quad \text{given} \quad \Delta^2 \geq \frac{E(X^2)}{E(X)^2}$$

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**Our contribution:**

$$\frac{\Delta}{\epsilon} \cdot \log^3 \left( \frac{B}{E(X)} \right)$$



# Our Approach

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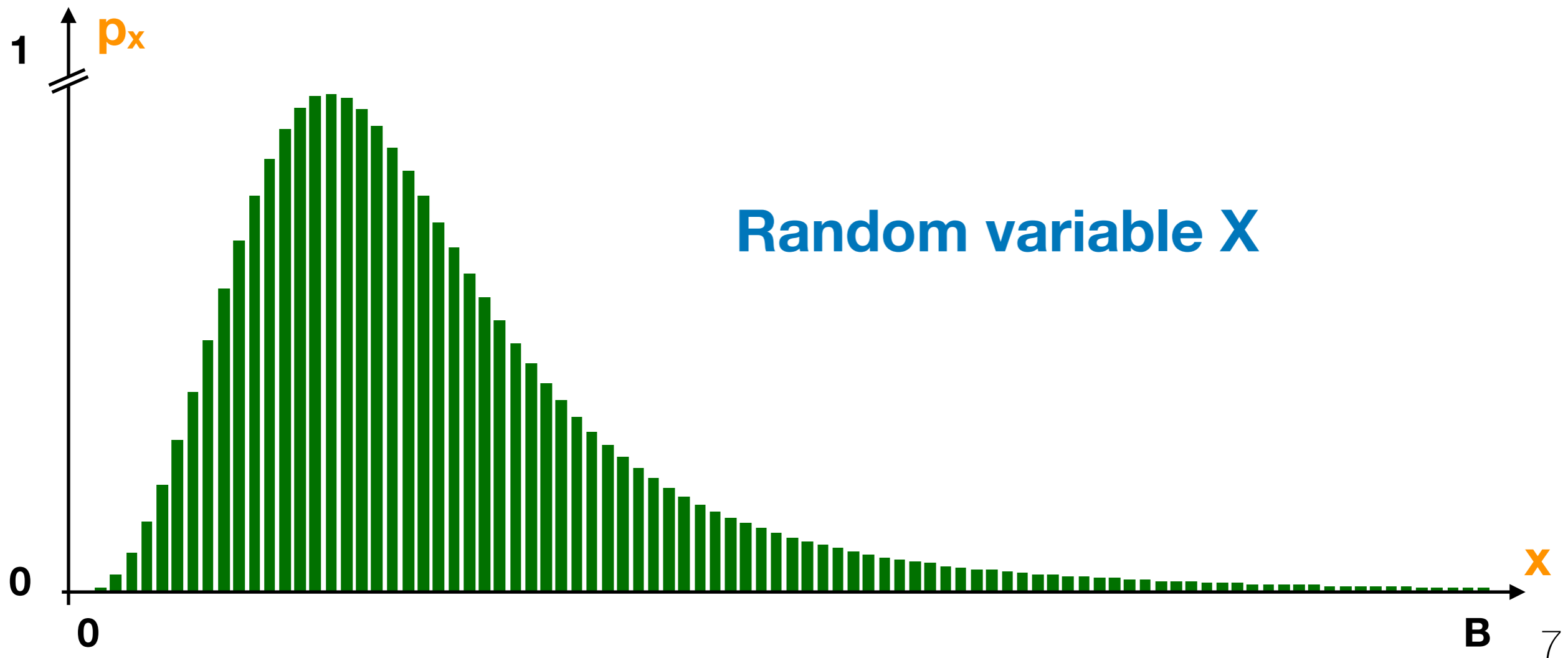
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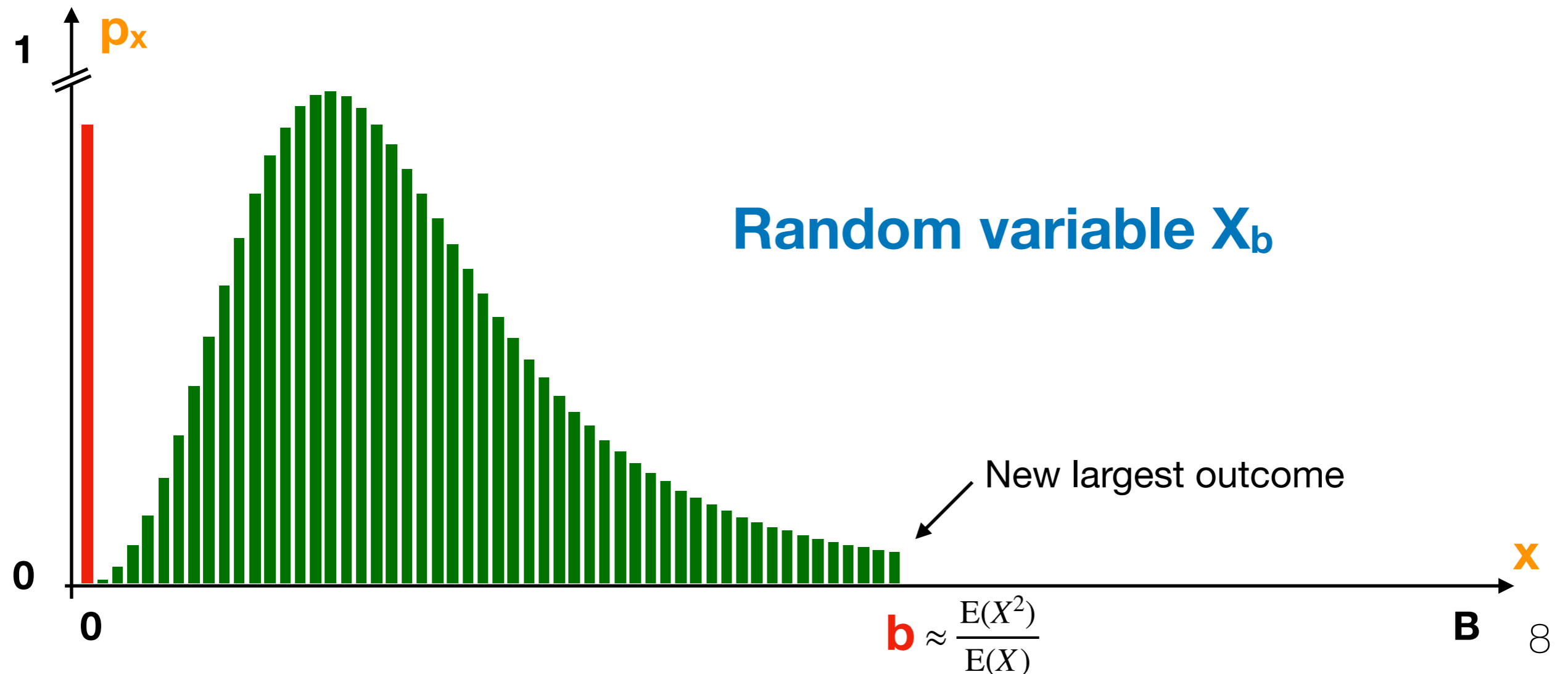
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Threshold	Input r.v.	Number of samples	Amplitude Estimation
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[Brassard et al.'02]

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# Applications

# Application 1: approximating graph parameters

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**Input:** graph  $G=(V,E)$  with  $n$  vertices,  $m$  edges,  $t$  triangles

**Query access:** unitaries  $O_{\text{deg}} |v\rangle |0\rangle = |v\rangle |\text{deg}(v)\rangle$  *(degree query)*

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$\rightarrow \tilde{\Theta} \left( \frac{\sqrt{n}}{t^{1/6}} + \frac{m^{3/4}}{\sqrt{t}} \right)$  quantum queries to **triangle** estimation

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[Goldreich, Ron'08] [Seshadhri'15]

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[Eden, Levi, Ron'15] [Eden, Levi, Ron, Seshadhri'17]

## Application 2: frequency moments in the streaming model

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**Initially:**  $\mathbf{x} = (0, \dots, 0)$  of **dimension**  $n$

**Input:** stream of updates  $\mathbf{x}_i \leftarrow \mathbf{x}_i + \delta$  to  $\mathbf{x}$

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**Result:**  $M = \tilde{O}\left(\frac{n^{1-2/k}}{P^2}\right)$  qubits of memory

(vs.  $M = \tilde{\Theta}\left(\frac{n^{1-2/k}}{P}\right)$  classical bits of memory)

[Monemizadeh, Woodruff'10]  
[Andoni, Krauthgamer, Onak'11]

**Conclusion**

The **mean** of a random variable  $X$  can be estimated with **multiplicative**

**error  $\epsilon$**  using  $\tilde{O}\left(\frac{\Delta}{\epsilon} \cdot \log^3\left(\frac{M_\Omega}{\mathbf{E}(X)}\right)\right)$  **quantum samples**, given  $\Delta^2 \geq \frac{\mathbf{E}(X^2)}{\mathbf{E}(X)^2}$ .

**Lower bound:**  $\Omega\left(\frac{\Delta - 1}{\epsilon}\right)$  quantum samples

**or**  $\Omega\left(\frac{\Delta^2 - 1}{\epsilon^2}\right)$  copies of the state  $S_X|0\rangle = \sum_{x \in \Omega} \sqrt{p_x} |\psi_x\rangle |x\rangle$

### Open questions:

- Can we improve the complexity to  $O(\Delta/\epsilon)$  exactly?
- Lower bound for the Frequency Moments estimation problem?
- Other applications ?

**arXiv: 1807.06456**